CSI 33 Midterm Exam In-class Practice

Part 1 Answer True/False and Multiple Choice questions

1. Which of the following is a $\Theta(n)$ operation?

- (a) Sorting a list with Selection sort
- (b) Finding the ith item in a Python list.
- (c) Re-assigning the element at the end of a Python list.
- (d) Deleting an item from the middle of a Python list.

Answer: (d)

Explanation: Selection Sort is $\Theta(n^2)$ operation;

random access in Python list (finding ith item, myList[i]) is $\Theta(1)$ operation; re-assigning an element at the end of a Python list (myList[len(myList)-1] = ...) is also $\Theta(1)$ operation;

When a middle element is deleted, the "right half" of the values, about $\frac{n}{2}$ of them, must be shifted one

space to the left, which gives $\Theta(n)$ running time.

- 2. Which of the following is **not** true of Python dictionaries?
 - (a) They are implemented as hash tables.
 - (b) Lookup is very efficient.

(c) Values must be immutable.

(d) All of the above are true.

Answer: (c)

Explanation: Indeed, Python dictionaries are implemented as hash tables, the lookup, insertion, deletion are all $\Theta(n)$ operations. Keys must be immutable (as this is the way to "access" the associated value with it). Values can me mutable.

3. How many iterations will the while loop of the *Binary Search* do when searching for 21 in the sequence [1, 5, 12, 14, 17, 21, 28]? Use the Binary Search algorithm I presented in class.

(a) 5 (b) 4 (c) 3 (d) 2

Answer: (d)

Explanation: look at the algorithm of the *Binary Search* for the key information:

1) the middle value accessed by the index $\left\lfloor \frac{high+low}{2} \right\rfloor$, where low = 0 and high = len(myList) - 1

initially

2) the while loop stops as soon as *low* > *high*

3) when the middle element checked for equality with the target value:

- if it is equal, then the index is returned, and

- if not, left half (stepping one left for the high index) or the right half (stepping one to the right for the low index) is "chosen"

21 is present in the sequence, therefore, the exit condition from the loop will be the location of this element.

We will begin by selecting index $\left\lfloor \frac{6+0}{2} \right\rfloor = 3$, the value at the 3rd position is 14, not 21.

Since 21 is greater than 14, the *low* index is adjusted to low = 3+1 = 4.

2nd iteration of the while loop: the "middle index" is

" is $\left\lfloor \frac{6+4}{2} \right\rfloor = 5$, the value at the 5th position is 21. It is the

target value, therefore, the position 5 is returned and the while loop is terminated. The Binary Search algorithm performed 2 iterations of the while loop.

Here is what you can present as an explanation, if asked for:



Part 2. Answer short-answer questions

1. Consider the following code fragment:

from ListNode import *

z = ListNode(34) y = ListNode(25,z) x = ListNode(12,y) t = ListNode(20,y)

What will be produced by this code fragment (draw a pictorial representation)?

For your reference, the definition of the ListNode class:

... skipped

Answer: graphical



2. Give a theta analysis of the time efficiency of the following code fragment. Provide explanations.

More details: note that initially the list is empty, and the next element is inserted into the "head" position (0th position), so we cannot claim that at each iteration the list has n elements and all of them are shifted 1 space to the right.

So it is better to see what is going on at each iteration from the very beginning:

1st iteration, the value of n is inserted into an empty list : 1 step

2nd iteration: the value of myList[0] is shifted one space to the right, and the n-3 is placed into the 0th position: 2 steps

3rd iteration: the values of myList[1] and the myList[0] are shifted one position to the right, and the value of n-6 is inserted into the 0th position: 3 steps

4th iteration: the values of myList[2], myList[1], and myList[0] are shifted to the right, and the value of n-9 is inserted into the 0th position: 4 steps

•••

the loop will stop when $n-3k \le 1...$ so there will be about $\frac{n}{2}$ iterations:

1 step + 2 steps + 3 steps + 4 steps + ... + $\frac{n}{3}$ steps = arithmetic sequence = $\frac{\left(1 + \frac{n}{3}\right)\left(\frac{n}{3}\right)}{2} = \dots = \frac{n}{6} + \frac{n^2}{6}$

Hence T(n) = $4 + 2 \cdot \left(\frac{n}{6} + \frac{n^2}{6}\right) = 4 + \frac{n}{3} + \frac{n^2}{3} = \Theta(n^2)$

Answer: $T(n) = \Theta(n^2)$

3. Give pictorial representation of the Python's memory during execution of the code given below. Show the result of print statements.



['a','b', 25] Hello 25

Part III.

Here is the running time for each of four requested operations:

operations	(a) an unordered Python list	(b) a sorted Python list	(c) a Python dictionary (elements of the set are keys, None or True are values)	
add	$\Theta(n)$ we need to check if the element is already in the set, and since the list is unordered, we will have to apply linear search	$\Theta(\log n)$ since the elements are ordered, we need to find a position to insert the new record, search can be done with <i>log n</i> time (recall binary search on sorted arrays), then append operation on average takes $\Theta(1)$ time	$\Theta(1)$ almost all basic operations on dictionaries are $\Theta(1)$, since hash tables with hashing function are used.	
remove	$\Theta(n)$ operations of insertion and deletion are $\Theta(n)$ for Python's lists + we need to find an element, and shift all the ones to the right of it one space to the left	$\Theta(n)$ first we will need to locate the element with the given name ($\Theta(\log n)$ operation), then we will need to delete is ($\Theta(n)$ operation on Pyton's lists), hence the result is $\Theta(n)$	$\Theta(1)$ using hashing function the record will be accessed in constant time, and deleted	
clear	$\Theta(1)$ if we reassign the data attribute to empty list, e.g. data=[]	$\Theta(1)$ if we reassign the data attribute to empty list, e.g. data=[]	Θ(n) either deleting all elements or changing the values of the keys to None	
contains	$\Theta(n)$ since the list is unordered, we will have to apply linear search	$\Theta(\log n)$ search can be done with <i>log</i> <i>n</i> time (recall binary search on sorted arrays)	Θ(1) one of the basic operations of Python dictionaries	
intersection	$\Theta(nm)$ set1 = n and set2 = m Every element from the set1 will be "searched for" in set2	Θ(n log m) set1 = n and set2 = m Every element from the set1 will be "searched for" in set2, w can use binary search	If n is number of elements in set1 and m is the number of elements of set2, and n < m, then we say the asymptotic running time is $\Theta(n)$ set1 = n and set2 = m, and n <= m	
intersection the idea is to grab the smaller size set and check if its elements are present in the other, if present, then the element is added to the new set. Worse case scenario – all the elements are to be added				
union	$\Theta(nm)$ We will start by adding all the elements of one set to the new set; then we will be	$\Theta(n+m)$ Two sets are ordered, so we compare the two first elements of the sets, and	⊖(n+m) All the keys from one dictionary are added to the new dictionary, then each	

	grabbing one by one elements	grab the smallest – add it to	key from the second
	from the second set, checking	the new set; then we	dictionary is checked for
	if its presence in the new set	compare the "next"	presence in the new
	– if it is not present, we will	front/first elements : take	dictionary: if is already
	add it, if it is – we will move	the smallest and add it to	there, don't add it,
	on to the next.	the new set, etc.	otherwise, add it.
difference	Θ(nm)	Θ(nm)	$\Theta(n)$
	Set1-Set2: grab an element	Set1-Set2: grab an element	Set1-Set2: grab an element
	from Set1 (n elements): if it is	from Set1 (n elements): if it	from Set1 (n elements): if it
	not in Set2 (m elements to	is not in Set2 (m elements	is not in Set2 ($\Theta(1)$
	compare to), add it to the new	to compare to), add it to the	operation), add it to the new
	set, and so forth.	new set, and so forth.	set, and so forth.

About operations and their cost on Python's list see pages 99 About operations and their cost on dictionaries see pages 94-95